

Functional Data Structures

Exercise Sheet 13

The following are old exam questions!

Exercise 13.1 Amortized Complexity

A “stack with multipop” is a list with the following two interface functions:

fun *push* :: “ $'a \Rightarrow 'a \text{ list} \Rightarrow 'a \text{ list}$ ” **where**
“*push* $x \ xs = x \# \ xs$ ”

fun *pop* :: “ $\text{nat} \Rightarrow 'a \text{ list} \Rightarrow 'a \text{ list}$ ” **where**
“*pop* $n \ xs = \text{drop } n \ xs$ ”

You may assume

definition *T_push* :: “ $'a \Rightarrow 'a \text{ list} \Rightarrow \text{nat}$ ” **where**
“*T_push* $x \ xs = 1$ ”

definition *T_pop* :: “ $\text{nat} \Rightarrow 'a \text{ list} \Rightarrow \text{nat}$ ” **where**
“*T_pop* $n \ xs = \min n \ (\text{length } xs)$ ”

Use the potential method to show that the amortized complexity of *push* and *pop* is constant.

If you need any properties of the auxiliary functions *length*, *drop* and *min*, you should state them but you do not need to prove them.

Exercise 13.2 Converting List for Balanced Insert

Recall the standard insertion function for unbalanced binary search trees.

fun *insert* :: “ $'a::\text{linorder} \Rightarrow 'a \text{ tree} \Rightarrow 'a \text{ tree}$ ” **where**
“*insert* $x \ \text{Leaf} = \text{Node } \text{Leaf } x \ \text{Leaf}$ ” |
“*insert* $x \ (\text{Node } l \ a \ r) =$
 (*case cmp* $x \ a$ of
 LT $\Rightarrow \text{Node } (\text{insert } x \ l) \ a \ r$ |
 EQ $\Rightarrow \text{Node } l \ a \ r$ |
 GT $\Rightarrow \text{Node } l \ a \ (\text{insert } x \ r)$)”

We define the function *from_list*, which inserts the elements of a list into an initially empty search tree:

definition *from_list* :: "*a::linorder list* \Rightarrow '*a tree*" **where**
"*from_list l = fold insert l Leaf*"

Your task is to specify a function *preprocess*::'*a*, that preprocesses the list such that the resulting tree is almost complete.

You may assume that the list is sorted, distinct, and has exactly $2^k - 1$ elements for some *k*. That is, your *preprocess* function must satisfy:

fun *preprocess* :: "*a list* \Rightarrow '*a list*"

lemma

assumes "*sorted l*"
and "*distinct l*"
and "*length l = 2^k - 1*"
shows "*set (preprocess l) = set l*" **and** "*acomplete (from_list (preprocess l))*"

Note: No proofs required, only a specification of the *preprocess* function!