

Functional Data Structures

Exercise Sheet 2

Exercise 2.1 Fold function

The fold function is a very generic function, that can be used to express multiple other interesting functions over lists.

Have a look at Isabelle/HOL's standard function *fold*.

thm *fold.simps*

Write a function to compute the sum of the elements of a list. Define two versions, one direct recursive definition, and one using fold. Show that both are equal.

fun *list_sum* :: "*nat list* \Rightarrow *nat*"

definition *list_sum'* :: "*nat list* \Rightarrow *nat*"

To use your definition in a proof, you need to use the theorem *list_sum'_def* explicitly.

lemma "*list_sum xs = list_sum' xs*"

Exercise 2.2 Folding over Trees

Define a datatype for binary trees that store data only at leafs.

datatype *'a ltree* =

Define a function that returns the list of elements resulting from an in-order traversal of the tree.

fun *inorder* :: "*'a ltree* \Rightarrow *'a list*"

In order to fold over the elements of a tree, we could use *fold f (inorder t) s*.

Define a function *fold_ltree* that is recursive on the structure of the tree, and that returns the same result as *fold f (inorder t) s*.

fun *fold_ltree* :: "*('a* \Rightarrow *'s* \Rightarrow *'s*) \Rightarrow *'a ltree* \Rightarrow *'s* \Rightarrow *'s*"

lemma "*fold f (inorder t) s = fold_ltree f t s*"

Define a function *mirror* that reverses the order of the leafs, i.e. that satisfies the following specification:

lemma “*inorder (mirror t) = rev (inorder t)*”

Exercise 2.3 Shuffle Product

A shuffle of two lists, *xs* and *ys*, is a list that contains exactly the elements of *xs* and *ys* s.t. every two elements $x \in xs$ (resp. *ys*) and $x' \in xs$ (resp. *ys*) occur in the shuffle in the same order they do in *xs* (resp. *ys*).

Define a function *shuffles* that returns a list of all shuffles of two given lists

fun *shuffles* :: “*a list* \Rightarrow *'a list* \Rightarrow *'a list list*”

Show that the length of any shuffle of two lists is the sum of the length of the original lists.

lemma “*zs* \in *set (shuffles xs ys)* \implies *length zs = length xs + length ys*”

Homework 2.1 Association Lists

Submission until Thursday, May 02, 23:59pm.

An association list is a list of pairs. An entry (k,v) means that key *k* is associated to value *v*.

For an association list *xs*, the *collect k xs* operation returns a list of all values associated to key *k*, in the order stored in the list. Specify the function *collect* by a set of recursion equations:

fun *collect* :: “*a* \Rightarrow (*'a* \times *'b*) *list* \Rightarrow *'b list*”

Test cases

definition *ctest* :: “(*int* * *int*) *list*” **where** “*ctest* \equiv [
 (2,3),(2,5),(2,7),(2,9),
 (3,2),(3,4),(3,5),(3,7),(3,8),
 (4,3),(4,5),(4,7),(4,9),
 (5,2),(5,3),(5,4),(5,6),(5,7),(5,8),(5,9),
 (6,5),(6,7),
 (7,2),(7,3),(7,4),(7,5),(7,6),(7,8),(7,9),
 (8,3),(8,5),(8,7),(8,9),
 (9,2),(9,4),(9,5),(9,7),(9,8)
]”

value “*collect 3 ctest = [2,4,5,7,8]*”

value “*collect 1 ctest = []*”

An experienced functional programmer might also write this function as

```
map snd (filter (λkv. fst kv = x) ys)
```

Show that this specifies the same function:

lemma *collect_alt*: “*collect x ys = map snd (filter (λkv. fst kv = x) ys)*”

When the lists get bigger, efficiency might be a concern. To avoid stack overflows, you might want to specify a tail-recursive version of *collect*. The first parameter is the accumulator, that accumulates the elements to be returned, and is returned at the end.

Note: To avoid appending to the accumulator, we accumulate the elements in reverse order, and reverse the accumulator at the end.

Complete the second equation!

```
fun collect_tr :: “'a list ⇒ 'b ⇒ ('b × 'a) list ⇒ 'a list” where  
  “collect_tr acc x [] = rev acc”
```

Show correctness of your tail-recursive version. Hint: Generalization!

lemma *collect_tr_collect*: “collect_tr [] x ys = collect x ys”

Homework 2.2 Perfectly Balanced Trees

Submission until Thursday, May 02, 23:59pm.

Recall the tree datatype *'a ltree* from the tutorial. Define functions to return the height (A leaf has height 0) and the number of leaves:

```
fun lheight :: “'a ltree ⇒ nat”  
fun num_leafs :: “'a ltree ⇒ nat”
```

A tree is perfectly balanced iff, for each node, the left and right subtree have the same height. Specify a function to check that a tree is perfectly balanced.

```
fun perfect :: “'a ltree ⇒ bool”
```

Show that, for a perfectly balanced tree with height *h* and number of leafs *l*, we have $l = 2^h$:

lemma *perfect_num_leafs_height*: “perfect t ⇒ num_leafs t = 2^{height t}”

Homework 2.3 Shuffles

Submission until Thursday, May 02, 23:59pm.

In the tutorial you have defined the function *shuffles*. Show that the elements in a shuffle are exactly the elements of the input lists.

lemma *set_shuffles*: “zs ∈ set (shuffles xs ys) ⇒ set zs = set xs ∪ set ys”