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Lambda Calculus
Winter Term 2017/18
Exercise Sheet 10

Exercise 1 (Recursive `let`)

Recursive `let` expressions are one way (besides Y-combinators) to add recursion to $\lambda^{\to}$.

$$t ::= x \mid (t_1\ t_2) \mid (\lambda x.\ t) \mid \texttt{letrec}\ x = t_1\ \texttt{in}\ t_2$$

a) Modify the standard typing rule for `let` to create a suitable rule for `letrec`.

b) Considering *type inference*, what is the problematic property of this rule compared to the rule for `let`?

c) Give a derivation tree for the following statement, and so determine the type τ:

$$[] \vdash\ \texttt{letrec}\ x = \lambda y.\ x\ (x\ y)\ \texttt{in}\ x\ x : \tau$$

Exercise 2 (Type Inference in Haskell (2))

In this exercise, we will extend the implementation of the type inference algorithm from last exercise for the `let` and `letrec` constructs.

a) Have a look at the template on the website, which adds `let` and `letrec` to the term language.

b) Extend the type inference algorithm for `let` and `letrec`.

Homework 3 (Fixed-point combinator)

Let

$$\$ = \lambda abcdefghijklmnopqstuvwxyzr.r(thisisafixedpointcombinator)$$

and

$$€ = \$.$$

Show that € is a fixed-point combinator.

Homework 4 (`let`-Polymorphism)

Give a derivation tree for the following statement, and so determine the type τ:

$$[z : \tau_0] \vdash\ \texttt{let}\ x = \lambda yz.\ z\ y\ y\ \texttt{in}\ x\ (x\ z)\ : \tau$$

Homework 5 (Type Inference with Type Constructors)

We generalize $\rightarrow$ to type constructors. With type constructors, types are either elementary, a type variable or, constructed as $\Pi\ \tau_1 \ldots \tau_n$ where Π is a type constructor. Now $\rightarrow$ is just a type constructor that takes two arguments. Your task is to extend the Haskell inference algorithm towards type constructors.

a) Extend the type language with type constructors. Type constructors should take *lists* of type arguments.

b) Extend the type inference algorithm for type constructors. To specify the set of valid type constructors, we will just start with a non-empty environment that will pre-specify the type of some free variables that then act as data constructors.