# Basic Proof Theory
## Propositional Logic

(See the book by Troelstra and Schwichtenberg)

# Proof rules and proof systems

Proof systems are defined by (proof or inference) rules of the form

$$\frac{T_1 \quad \ldots \quad T_n}{T} \text{ rule-name}$$

where $T_1, \ldots T_n$ (premises) and $T$ (conclusion) are syntactic objects (eg formulas).

Intuitive reading: If $T_1, \ldots, T_n$ are provable, then $T$ is provable.

Degenerate case: If $n = 0$ the rule is called an axiom and the horizontal line is sometimes omitted.

If some $U$ is provable, we write $\vdash U$.

# Proof trees

Proofs (also: derivations) are drawn as trees of nested proof rules.

Example (Proof/derivation tree)

$$\dfrac{\dfrac{}{T_1} \quad \dfrac{\dfrac{}{U}}{T_2}}{S_1} \qquad \dfrac{\dfrac{}{T_3}}{S_2}$$
$$\dfrac{S_1 \qquad S_2}{R}$$

We sometimes omit the names of proof rules in a proof treeif they are obvious or for space reasons. *You* should always show them!

Every fragment

$$\frac{T_1 \quad \ldots \quad T_n}{T}$$

of a proof tree must be (an instance of) a proof rule.
All proofs must start with axioms.

The depth of a proof tree is the number of rules on the longest branch of the tree. Thus $\geq 1$

# Abbreviations

Until further notice:

$\bot$, $\neg$, $\wedge$, $\vee$, $\rightarrow$ are primitives.

$\top$ abbreviates $\neg\bot$

A possible simplification:

$\neg F$ abbreviates $F \rightarrow \bot$

We now consider three important proof systems:

- Sequent Calculus
- Natural Deduction
- Hilbert Systems