

LOGICS EXERCISE

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EXERCISE SHEET 11

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**Submission of Homework:** Before tutorial on June 29

**Exercise 11.1.** [Natural Deduction (Warmup)]

Prove by natural deduction:

1.  $(F \wedge G) \wedge H \rightarrow F \wedge (G \wedge H)$

2.  $(F \vee G) \vee H \rightarrow F \vee (G \vee H)$

**Exercise 11.2.** [Natural Deduction (Advanced)]

Prove by natural deduction:

1.  $\neg(F \wedge G) \rightarrow (\neg F \vee \neg G)$

2.  $((F \rightarrow G) \rightarrow F) \rightarrow F$

**Exercise 11.3.** [Alternative  $\wedge$ E rule]

Show how to transform a natural deduction proof that additionally uses the following rule to one that does not use the rule:

$$\frac{F \wedge G \quad \begin{array}{c} [F, G] \\ \vdots \\ H \end{array}}{H}$$

**Homework 11.1.** [Natural Deduction (Warmup)] (3 points)

Show:

$$\vdash_N (F \rightarrow G) \rightarrow (\neg G \rightarrow \neg F)$$

**Homework 11.2.** [Classical Reasoning (1)] (6 points)

We replace rule  $\perp$  of the calculus of natural deduction by either one of the following rules:

- $\frac{}{F \vee \neg F}$  (law of excluded middle)
- $\frac{\neg\neg F}{F}$  (double negation elimination)

Additionally, we add the rule  $\frac{\perp}{F}$  ( $\perp E$ ). Show that the calculus of natural deduction remains complete in both cases.

**Homework 11.3.** [Classical Reasoning (2)] (5 points)

Assume that the calculus of natural deduction is *augmented* with the two rules from the last exercise. Show:

- $\vdash_N (\neg G \rightarrow \neg F) \rightarrow (F \rightarrow G)$
- $\vdash_N (\neg F \rightarrow G) \rightarrow (F \vee G)$

*Hint:* You can also use the law of excluded middle in the following form:

$$\frac{\begin{array}{c} [F] \\ \vdots \\ G \end{array} \quad \begin{array}{c} [\neg F] \\ \vdots \\ G \end{array}}{G}$$

**Homework 11.4.** [Left-Sided Sequent Calculus] (6 points)

We want to study a modified sequent calculus where the right-hand side is always empty, i.e. where sequents are of the form  $\Gamma \Rightarrow$ . Give a set of rules for this calculus such that your calculus fulfills the following property and sketch a proof:

$$\Gamma \Rightarrow \Delta \text{ iff } \Gamma, \neg\Delta \Rightarrow$$

*Hint:* Use induction over the length of the derivations. You can skip the cases for  $\vee$  and  $\wedge$  and instead look at  $\rightarrow$ .