

# LOGICS EXERCISE

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EXERCISE SHEET 6

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**Submission of homework:** Before tutorial on 13.06.2017. You have to do the homework yourself; no teamwork allowed.

**Exercise 6.1. [Unused Bound Variables]**

For this exercise, we first define an alternative way to evaluate a formula in a structure, based on arithmetic and set operations.

$$\begin{aligned}\mathcal{A}(\neg F) &= 1 - \mathcal{A}(F) \\ \mathcal{A}(F \vee G) &= \max(\mathcal{A}(F), \mathcal{A}(G)) \\ \mathcal{A}(F \wedge G) &= \min(\mathcal{A}(F), \mathcal{A}(G)) \\ \mathcal{A}(\exists x F) &= \max\{\mathcal{A}[d/x](F) \mid d \in U_{\mathcal{A}}\} \\ \mathcal{A}(\forall x F) &= \min\{\mathcal{A}[d/x](F) \mid d \in U_{\mathcal{A}}\}\end{aligned}$$

The evaluation for predicates and terms remain unchanged.

Equipped with this definition, prove the equivalence  $\exists x F \equiv F$  where  $x$  does not occur in  $F$ .

*Hint:* Adapt the coincidence lemma for propositional logic (exercise 1.1) to predicate logic.

**Exercise 6.2. [Substitution Lemma]**

Consider the following statement: “If  $F \equiv F'$ , then  $F[t/x] \equiv F'[t/x]$ .” Proof or counterexample.

**Exercise 6.3. [Skolem Form]**

Convert the following formula into – in order – a rectified formula, RPF and Skolem form.

$$P(x) \wedge \forall x (Q(x) \wedge \forall x \exists y P(f(x, y)))$$

**Exercise 6.4. [Herbrand Models]**

Given the formula

$$F = \forall x \forall y (P(f(x), g(y)) \wedge \neg P(g(x), f(y)))$$

1. Specify a Herbrand model for  $F$ .
2. Specify a Herbrand structure suitable for  $F$ , which is not a model of  $F$ .

**Homework 6.1. [Skolem Form]**

(6 points)

Convert the following formulas into – in order – a rectified formula, RPF and Skolem form.

1.  $\forall x \exists y \forall z \exists w (\neg P(a, w) \vee Q(f(x), y))$

2.  $\forall z \exists y (P(x, g(y), z) \vee \neg \forall x Q(x))$

**Homework 6.2. [Invalid Herbrand Models]**

(8 points)

Recall the fundamental theorem from the lecture: “Let  $F$  be a closed formula in Skolem form. Then  $F$  is satisfiable iff it has a Herbrand model”.

Explain “what goes wrong” if the precondition is violated: when  $F$  is not closed or not in Skolem form. Describe both cases.

**Homework 6.3. [Proof of the Fundamental Theorem]**

(6 points)

Recall the proof of the fundamental theorem from the lecture. Give the omitted proof for the base case (slide 6,  $\mathcal{A}(G) = \mathcal{T}(G)$ ).