

Semantics of Programming Languages

Exercise Sheet 1

Before beginning to solve the exercises, open a new theory file named `Ex01.thy` and write the the following three lines at the top of this file.

```
theory Ex01
imports Main
begin
```

Exercise 1.1 Calculating with natural numbers and integers

Use the `value` command to turn Isabelle into a fancy calculator and evaluate the following natural number expressions:

```
"2 + (2::nat)" "(2::int) * (5 + 3)" "(3::int) * 4 - 2 * (7 + 1)"
"(3::nat) * 4 - 2 * (7 + 1)"
```

Can you explain the last result?

Exercise 1.2 Natural number laws

Formulate and prove the well-known laws of commutativity and associativity for addition of natural numbers.

Exercise 1.3 Counting elements of a list

Define a function which counts the number of occurrences of a particular element in a list.

```
fun count :: "'a list ⇒ 'a ⇒ nat"
```

Test your definition of `count` on some examples and prove that the results are indeed correct.

Prove the following inequality (and additional lemmas, if necessary) about the relation between `count` and `length`, the function returning the length of a list.

```
theorem "count xs x ≤ length xs"
```

Exercise 1.4 Adding elements to the end of a list

Define a function *snoc* that appends an element at the right end of a list. Do not use the existing append operator `@` for lists.

```
fun snoc :: "'a list ⇒ 'a ⇒ 'a list"
```

Convince yourself on some test cases that your definition of *snoc* behaves as expected, for example run:

```
value "snoc [] c"
```

Also prove that your test cases are indeed correct, for instance show:

```
lemma "snoc [] c = [c]"
```

Next define a function *reverse* that reverses the order of elements in a list. (Do not use the existing function *rev* from the library.) Hint: Define the reverse of $x \# xs$ using the *snoc* function.

```
fun reverse :: "'a list ⇒ 'a list"
```

Demonstrate that your definition is correct by running some test cases, and proving that those test cases are correct. For example:

```
value "reverse [a, b, c]"
```

```
lemma "reverse [a, b, c] = [c, b, a]"
```

Prove the following theorem. Hint: You need to find an additional lemma relating *reverse* and *snoc* to prove it.

```
theorem "reverse (reverse xs) = xs"
```

Homework 1 Gauss Sequence

Submission until Tuesday, October 22, 10:00am.

Homework must be sent by mail to lammich@in.tum.de. Send a *working* theory file named FirstnameLastname01.thy, that contains no *sorry*-commands.

The pen-and-paper part may be submitted by email or handed in at the tutorial.

In this homework assignment you will prove the formula for the well-known Gauss sequence:

$$\sum_{i=0}^n i = \frac{n(n+1)}{2}$$

Your first task is to use the **fun** command to define a function *sum* :: *nat* ⇒ *nat* in Isabelle, that computes the sum of the first *n* natural numbers.

```
fun sum :: "nat ⇒ nat"
```

You may wish to use the **value** command to check that your definition is correct. For example, the following command should evaluate to *True*.

value “*sum 3 = 1+2+3*”

Now prove the Gauss sequence formula by induction over n . First, write an informal proof by hand. Your proof should contain a base case for zero, where you show that $sum(0)$ equals 0 . Next you should have a case for successor: Fix an arbitrary m , assume the inductive hypothesis that $sum(m) = m*(m+1) \text{ div } 2$, and then show that $sum(Suc\ m) = (m+1)*(m+2) \text{ div } 2$.

Finally, prove the same property *formally* in Isabelle:

lemma “*sum n = n * (n+1) div 2*”