

## Semantics of Programming Languages

### Exercise Sheet 10

#### **Exercise 10.1** Using the VCG

Use the VCG to prove correct a multiplication and a square root program:

\*Note: With the template, use *unbundle COM/unbundle ACOM* before writing down a *com/acom* definition, otherwise parsing will be slow.

**definition MUL :: com where**

```
"MUL =
  "z'':=N 0;;
  "c'':=N 0;;
  WHILE (Less (V "c'') (V "y'')) DO (
    "z'':=Plus (V "z'') (V "x'');;
    "c'':=Plus (V "c'') (N 1))"
```

**theorem MUL\_partially\_correct:**

```
"{λs. 0 ≤ s "y" ∧ s=sorig}
  MUL
  {λs. s "z" = s "x" * s "y" ∧ (∀ v. v∉{"z","c"}) → s v = sorig v}"
```

**definition SQRT :: com where**

```
"SQRT =
  "r'':= N 0;;
  "s'':= N 1;;
  WHILE (Not (Less (V "x") (V "s'))) DO (
    "r'':= Plus (V "r") (N 1);;
    "s'':= Plus (V "s") (V "r");;
    "s'':= Plus (V "s") (V "r");;
    "s'':= Plus (V "s") (N 1)
  )"
```

**theorem SQRT\_partially\_correct:**

```
"{λs. s=sorig ∧ s "x" ≥ 0}
  SQRT
  {λs. (s "r")^2 ≤ s "x" ∧ s "x" < (s "r"+1)^2 ∧ (∀ v. v∉{"s","r"}) → s v = sorig v}"
```

## Exercise 10.2 Total Correctness

Prove total correctness of the multiplication and square root program.

Rotated rule for sequential composition:

```
lemmas Seq_bwd = Hoare_Total.Seq[rotated]
```

Prove the following syntax-directed conditional rule (for total correctness):

```
lemma IfT:
assumes " $\vdash_t \{P1\} c_1 \{Q\}$ "  

and " $\vdash_t \{P2\} c_2 \{Q\}$ "  

shows " $\vdash_t \{\lambda s. (bval b s \rightarrow P1 s) \wedge (\neg bval b s \rightarrow P2 s)\} \text{ IF } b \text{ THEN } c_1 \text{ ELSE } c_2 \{Q\}$ "
```

```
lemmas hoareT_rule[intro?] = Seq_bwd Hoare_Total.Assign Hoare_Total.Assign' IfT
```

**theorem MUL\_totally\_correct:**

```
" $\vdash_t \{\lambda s. 0 \leq s "y" \wedge s = sorig\}$   

MUL  

 $\{\lambda s. s "z" = s "x" * s "y" \wedge (\forall v. v \notin \{"z", "c"\} \rightarrow s v = sorig v)\}$ "
```

**theorem SQRT\_totally\_correct:**

```
" $\vdash_t \{\lambda s. s = sorig \wedge s "x" \geq 0\}$   

SQRT  

 $\{\lambda s. (s "r")^2 \leq s "x" \wedge s "x" < (s "r" + 1)^2 \wedge (\forall v. v \notin \{"s", "r"\} \rightarrow s v = sorig v)\}$ "
```

## Homework 10.1 Squaring (5 pts)

Submission until Monday, Jan 16, 23:59pm.

Consider the following program for squaring a non-negative integer:

```
z := 1;  

a := 0;  

WHILE 0 < n DO (  

    a := a + z;  

    z := z + 2;  

    n := n + (-1)  

)
```

Here is a picture illustrating the algorithm:

```
1 2 3 4  

2 2 3 4  

3 3 3 4  

4 4 4 4
```

Using the VCG, prove that this algorithm correctly squares non-negative integers:

**theorem** *SQUARE\_correct*:

$$\text{“} \vdash \{\lambda s. s \text{ “}n\text{”} \geq 0 \wedge s = s_0\} \text{ } \text{SQUARE} \{ \lambda s. s \text{ “}a\text{”} = s_0 \text{ “}n\text{”} * s_0 \text{ “}n\text{”}\} \text{”}$$

Hints:

- Your invariant should state something about all three variables,  $n$ ,  $a$ , and  $z$ .
- *algebra\_simps* can help to solve any remaining verification conditions.

## **Homework 10.2** Be Original!

*Submission until Monday, Jan 16, 23:59pm.*

We're giving you one more week to finalize and polish your developments!

*In total, this exercise will be worth 15 points, plus bonus points for nice submissions.*